IN THE SPECIFICATION:

Amend p. 5, third paragraph, ln 6 – 8 as follows:

FIGURES 19A-E together show [[is]] a schematic diagram of a copier multiplexed keyboard with remote keystroke monitoring and remote keystroke operation capabilities.

Amend p. 5, fourth paragraph, ln 9 - 10 as follows:

FIGURE 19[[A]] F is a multiplexed keyboard timing diagram.

Amend p. 5, fifth paragraph, ln 11 – 12 as follows:

FIGURES 20[[A]] a1 and a2 [[is]] are an overview menu selection chart for use with the central computer.

Amend p. 16, third paragraph as follows:

The process is shown in Figures 19A-E (all of which together form a single schematic drawing). A keyboard S1 – S32, is shown organized as a 4 X 8 matrix. The rows as strobed by signals ROW*1 – ROW*4 with returning sense columns COL0 – COL7 pulled to a logical 1 by resistors R1 – R8, respectively. The resultant column data is read via the copier control to obtain D0 – D7 through the buffer U1 for each occurrence of a row scan signal when the COL READ EN* signal is asserted. An example of this operation is found in the multiplexed timing diagram of FIGURE 19[[A]]F. When an operator presses a key, such as S10 (copier start) at some time T1, and the keyboard matrix ROW2* is asserted at time T3 during a scan cycle, the COL2 sense line will be driven to a logical 0. This corresponds to the data bit D1 when buffer U1 is read.

Amend p. 17, first paragraph as follows:

Because the keyboard scanning operating may be too brief for the translator 6 to capture, a set of four 8-bit latches, U6 – U9 corresponding to each matrix row, are provided to automatically capture the column sense signals. As each row is strobed by ROW1* - ROW4*, the sense column data is clocked into its respective latch. In this way, the translator CPU 22 can read each row latch by asserting signals RD ROW 1 DATA* - RD ROW 4 DATA* asynchronous to the actual copier scanning operation to obtain a current image of the switch matrix. This image is then evaluated for copier specific information before being transmitted to the MAC computer 16. In this timing diagram, FIGURE 19[[A]]F, translator 6 reads an $FF (hexadecimal) from the latch U8, between T4 and T5, to obtain the image of switches S17 – S24. The image indicates that no keys are currently pressed within row 3.

Amend p. 18, first paragraph as follows:

In the keyboard timing diagram FIGURE 19[[A]]F, an example is shown for the assertion of switch S10. The translator CPU 22 writes a $FD into latch U3. When ROW2* is asserted, COL1 will be driven to a logical 0 which corresponds to switch S10 being pressed. The image that the copier control computer will receive on D1 when U1 is read will appear as if S10 has been pressed by an operator. After a predetermined period of time, the latches U2 – U5 are cleared to remove the simulated keystroke.

Amend p. 23, last paragraph as follows:

The data collection computer 16 can be an IBM compatible personal computer consisting of a monitor, keyboard, CPU, floppy drive, hard disk drive, and 640K of Random Access Memory running DOS 3.3. The data collection computer 16 assembles the status information into various display formats. Some of the user features are displayed in the Menu Selection Tree (FIGURES 20[A]a1 – 20[E]e). These features enable a database of information on copiers by manufacturer, model, options, location, facilities, etc. to be built. The database would then be merged with status information to present a current representation of status of all copiers 2 on the monitoring network. Copiers 2 with operational problems are easily identified and service requests made and tracked in like manner.